

§ 3. Algorithm of Resolution

Now we are ready to prove

$$\text{Ord II in dim} \leq n-1 \xrightarrow{\Pi_1} \text{Ord I in dim } n \xrightarrow{\Pi_2} \text{Ord II in dim } n.$$

Π_1 : We start from (X, I, E) $E = (\overset{\text{div smooth}}{E^1}, \dots, E^s)$ $\text{max ord } I \leq m$
 If $\text{max ord} < m$, the process is trivial. We assume $\text{max ord } I = m$

Step 1: Construct $\Pi_1: (X_{r_1}, I_{r_1}, E_{r_1}) \rightarrow \dots \rightarrow (X, I, E)$ s.t.

$$\text{Supp } \Pi_{1,x}^{-1} E \cap \text{Supp } (I_{r_1}, m) = \emptyset.$$

Consider the equivalent ideal $C(I)$ now

1.1. Let Z_0 be the union of all irr comp of E' contained in $\text{Supp } (I, m)$.

Blow up Z_0 , we have (for irr comp E^{1k} in E' for example) $\Pi_0: X_i \rightarrow X$

$$\text{max ord}_{E^{1k}} I = m \Rightarrow \text{max}_{\Pi_{0,x}^{-1} E^{1k}} \Pi_{0,x}^{-1} I \leq m - m = 0.$$

$\Rightarrow \Pi_{0,x}^{-1} E^{1k}$ and $\text{supp } (I_{r_1}, m)$ are disjoint.

1.2. Now, set $S = E^1$, $E_S = (E - E^1)|_S = (\mathcal{O}_S/E^1|_S, \dots, E^s|_S)$

consider $(S, I|_S, m, E|_S)$ apply Ord II in dim $\leq n-1$, we get

$$\mathcal{B}\mathcal{O}_1: (S_{E^1}, I_{E^1}|_{S_{E^1}}, m, E_{S_{E^1}}) \xrightarrow{\text{lift}} (X_{r_{E^1}}, I_{r_{E^1}}, m, E_S) \rightarrow \dots \rightarrow (\dots)$$

$$\text{s.t. } \text{supp } (I_{r_{E^1}}, m) \cap \Pi_{E^1,x}^{-1} S = \emptyset$$

$$\stackrel{\Pi_{E^1,x}^{-1} E^1}{\cong}$$

Inductively, we have $\Pi_1: (\dots) \rightarrow \dots (\dots)$ s.t. $\text{Supp } \Pi_{1,x}^{-1} E \cap \text{Supp } (I_{r_1}, m) = \emptyset$

Functionality follows from previous remark.

Step 2: Start from $J = H(C(I_{r_1}))$, $Y = X_{r_1}$, $F = (\Pi_{r_1}^{-1} E, \overset{\text{str}_1}{E^{s+1}}, \dots, E^{\text{str}_1})$
 (Y, J, m, F)

$\forall y \in \text{Supp } (J, m)$, $\exists U_y, h \in T(J)$ Since in step 1, all blow-up is snc, we can take h s.t. H_h snc w.r.t F ($\Pi_{r_1}^{-1} E$ away from $\text{supp } (I, m)$). \times

locally consider $(H_h, J|_{H_h}, m, F|_{H_h})$

Apply Ord II in dim $< n$ and lift it and globalized it to

$$\Pi_2: (X_{r_2}, I_{r_2}, E_{r_2}) \rightarrow (X_{r_1}, I_{r_1}, E_{r_1})$$

$$\text{s.t. } \text{Supp } (I_{r_2}, m) = \emptyset!$$

(Note, all blow up seq is also for $C(I)$ and I , $\text{TH}(C(I)) = C(I)$)

Functionality follows from previous rem.

$\star \uparrow \checkmark$
□

T2: We start with a marked triple (X, I, m, E) $E = (E^1, \dots, E^s)$.

Step 0:

We may write $I = N(I)M(I)$, $M(I) = \mathcal{O}_X(-\sum c_i E_i)$ and $\text{Supp } N(I)$ does not contain any of E^i !

Rem: if $E = \emptyset$, $M(I) = \emptyset$ and $I = N(I)$.

Step 1: Write $(N(I), 1) + (I, m) = (J, s)$ here s a number.

Write $m_J = \max \text{ord } J$, Run $\text{Ord } I$ to (X, J, E) with m_J .
we get (X, J, E_1) s.t. $\max \text{ord } J_1 = m_{J_1} < m_J$.

(Note, $\text{Supp } (J, m_J)_k \subseteq \text{Supp } (J, s)_k = \text{Supp } (N(I), 1)_k \cap \text{Supp } (I, m)_k$)

Inductively we get

(X_r, J_r, E_r) s.t. $\max \text{ord } J_r < s$.

this implies $\phi = \text{Supp } (J, s)_r = \text{Supp } (N(I), 1)_r \cap \text{Supp } (I, m)_r$
 $\text{Supp } (N(I)_r)$.

To sum up, we have $(X_r, I_r, m, E_r) \rightarrow \dots \rightarrow (X, I, m, E)$
s.t. $N(I)_r \cap \text{Supp } (I, m)_r = \emptyset$.

Rem: $N(I)_r$ $N(I_r)$ differs by some exceptional comps E^k ($k > s$)
 $M(I)_r$ $M(I_r)$ and is contained in $M(I_r)$.

Step 2: $I = M(I) = \mathcal{O}_X(-\sum a_j E^j)$ $E = (E^1, E^2, \dots, E^s, \dots)$.

2.1 $\text{Sub } \{E^1, E^2, \dots, E^s, \dots\}$ has a lexicographic order.

(x, \dots).

2.2. $\forall x \in X$, set $p(x) = (\{E^{j_1}, \dots, E^{j_k}\})$ the maximal subset (in above order) satisfying

(1) E^{j_i} pass $x \quad \forall 1 \leq i \leq k$.

(2) $\sum_{i=1}^k a_{j_i} \geq m$ (3) $\hat{a}_{j_1} + \dots + \hat{a}_{j_i} + \dots + a_{j_k} < m$.

$D_{p(x)} = \bigcap_{i=1}^k E^{j_i}$, and it is the locus that is a maximal component of

$\text{Supp } (I, m)$.

$v = (\max \text{ord } I, \text{ member of maximal comp of } \text{supp } (I, m) \text{ attain } \max \text{ord})$
 $= (m, n)$

Each time, we blow up $D_p(x)$, $D_p(x) = \{X_{j_1} = \dots = X_{j_k} = 0\}$.

$$\forall x \in D_p(x), \quad \text{since } \sum_{i=1}^k a_{j_i} - m < a_{j_i} \quad \forall 1 \leq i \leq k,$$

$$I_x = \prod_t x_t^{a_t}$$

$$\text{and } \pi_x^{-1} I_x = \prod_{\substack{t+j_i \\ \text{in coord}}} x_t^{a_t} \cdot x_{j_i}^{\sum a_{j_i} - m < a_{j_i}}$$

$$\Rightarrow \text{ord}_y \pi_x^{-1}(I) < \text{ord}_x I \quad \forall \pi(y) = x,$$

$\Rightarrow v$ decrease strictly in the Lexicographic order \Rightarrow the procedure terminates with $v = (m, \ast) \Rightarrow \max \text{ord } I_r < m$ eventually.

Functionality follows from the sm invariant property of $p(x)$ and v . □.

OCE: Commute resp to closed embedding or $\text{Ord II } m=1$ $E = \emptyset$

$$\tau: S \rightarrow X$$

Now, $\tau \times I_S$ in X we have locally $\tau \times I_S$ is adding some smooth element $\{U\}$.

$$\text{Ord}(\tau \times I_S, 1) = 1 \quad \tau(\tau \times I_S) = \tau \times I_S \leftarrow \text{sm element.}$$

Recall T_1 : We restrict to sm hypers in T to do induction,

so for this case, we just restrict $\tau \times I_S$ to $\{U\}$ and exactly get I , the procedure commute with closed embedding. □